# Global Alignment of Molecular Sequences via Ancestral State Reconstruction 

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# Capacity of noisy channels or Trace Reconstruction on a star 



- Choose $k$ random bits $x_{1} \ldots x_{k}$
$\square$ N - some noisy channel
$\square$ Goal: Given many applications of $N\left(x_{1} \ldots x_{k}\right)$ reconstruct $y_{1} \ldots y_{k}$ s.t.
$\operatorname{Pr}\left(x_{i}=y_{i}\right) \geq 0.99$
- How many channel uses do we need?


## Number of channel uses



- N applies i.i.d substitutions: constant number of uses (bit wise majority)
- N applies i.i.d. deletions, with constant probability - poly(k) uses [HMPW08]
$\square$ Both insertions and deletions, more general channels, subconstant probabilities - many open questions


## Trace reconstruction on a tree

- A recursive variant of trace reconstruction on a star
- On each edge, there is a probability for insertions, deletions and substitutions
- We are interested at
k sites in the root $\mathrm{x}_{1} \ldots \mathrm{x}_{\mathrm{k}}$
 a constant expansion ratio d


## Motivation

- Study of more general noisy channels
- Phylogenetic reconstruction

- Statistical Physics


## Main result

- Consider a d-ary tree, and a channel N :
- $N$ applies i.i.d substitutions with probability $\mathrm{p}_{\mathrm{s}^{\prime}}$ s.t.

$$
\left(1-2 p_{s}\right)^{2}>O\left(\frac{\log d}{d}\right)
$$

- $N$ applies i.i.d insertions with probability at most $O\left(1 / k^{2 / 3} h\right)$
- $N$ applies i.i.d insertions with probability at most $O\left(1 / k^{2 / 3} h\right)$
- Then one can "reconstruct" $x_{1} \ldots x_{k}$ from the leaves of the tree: Find $y_{1} \ldots y_{k}$ s.t. $\operatorname{Pr}\left(x_{i}=y_{i}\right)>0.99$
- Some lower bounds:
- Maximum substitution probability (without indels)

$$
\left(1-2 p_{s}\right)^{2}>1 / d
$$

- $1 / \mathrm{h}$ dependency corresponds to a constant fraction of deletions in the star case


## Recursive reconstruction (Mos98)

- Reconstruct the tree layer by layer
- Given d input vertices, reconstruct their father
- Continue recursively until the root is reconstructed
- Challenges:
- Sometimes the reconstruction fails
- Even when the reconstruction succeeds, and the children are perfect, the father is reconstructed up to some noise


## Reconstructing the father from the children

## $\square$ Each Child is divided into anchors and islands



## Reconstruction cntd.


$\square$ Align the children according to the anchors
$\square$ Do a place-wise majority on the children

## Where can we succeed?

- We can not reconstruct all the vertices correctly
- Suppose the first bit gets deleted going from the father to all the children
- Call a vertex v good if all three hold:
- There are no indel operations in the anchor, when going from $v$ to its children
- In each island of $v$, there is at most one indel operation
- It has at least d-2 good children
- The algorithm reconstructs all good vertices


## Correctness of the reconstruction

- Main Result follows from two theorems
- Thm 1: With high probability, the root is good
- Thm 2: The algorithm reconstructs all good vertices correctly


## Thm 1 - w.h.p the root is good

- Proof sketch: Show by induction on the height that most vertices are good.
- When is a vertex not good?
- When there are indels in the anchors: Improbable event, as anchors are short
- When there are two deletions in the same island: Improbable event, the islands have length $\mathrm{k}^{1 / 3}$, and the indel probability is $1 / \mathrm{k}^{2 / 3}$
- When two children are not good: improbable event by induction hypothesis
- Probability that the root is good $>0.99$


## Thm 2: reconstruction of good vertices

- Proof is by induction.
$\square$ Suppose v is good. All good children (> d-1) reconstructed "correctly"
$\square$ Reconstructed children + No indels in anchors $\rightarrow$ Alignment of the anchors is "correct"
- Correct alignment + each island suffers at most one deletion $\rightarrow y_{i}$ is a majority of $d$ values, such that d-2 of them are his true descendents
- Given that the majority is on the right descendents, we do not need to worry about the indels. Thm 2 holds because majority does error correction.


## Open questions

- Adversarial root
- Can we use these techniques to say something about the star case?
$\square$ Improving the parameters. In particular, a weaker definition of reconstruction, with higher deletion probabilities.
$\square$ Can we do something even without the tree?
- Follow up work shows how to reconstruct the topology of the tree (ABH'09)

Thank you

